

Inverting the signature of a path

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The signature of a path

Euclidan coordinates for \mathbf{R}^d : (e_1, \dots, e_d) .

A path $\gamma = (\gamma^1, \dots, \gamma^d) : [0, 1] \rightarrow \mathbf{R}^d$ continuously differentiable.

For word $w = e_{i_1} \cdots e_{i_n}$, define

$$C_\gamma(w) = \int_{0 < u_1 < \dots < u_n < 1} d\gamma_{u_1}^{i_1} \cdots d\gamma_{u_n}^{i_n}.$$

The signature of γ is the collection of all $C(w)$'s, denoted by $\text{Sig}(\gamma)$:

$$\text{Sig}(\gamma) := \sum_{n \geq 0} \underbrace{\sum_{w: |w|=n} C_\gamma(w) \cdot w}_{\text{Sig}^{(n)}(\gamma)}.$$

Independent of parametrisation. It captures the ordered evolution along the path through the order of the letters.

Examples

Consider \mathbf{R}^2 with standard basis $(e_1, e_2) = (x, y)$.

- ① Path movement: $(0, 0) \rightarrow (a, 0) \rightarrow (a, b)$.

$$ax * by \mapsto \exp(ax) \exp(by) .$$

- ② Path movement: $(0, 0) \rightarrow (0, b) \rightarrow (a, b)$.

$$by * ax \mapsto \exp(by) \exp(ax) .$$

- ③ Straightline segment from $(0, 0)$ to (a, b) .

$$ax + by \mapsto \exp(ax + by) .$$

They have the same $\text{Sig}^{(1)}$, but different $\text{Sig}^{(2)}$.

Another example

For every n , we give two paths α and β such that $\text{Sig}^{(k)}(\alpha) = \text{Sig}^{(k)}(\beta)$ for every $k \leq n$.

Consider dimension two. Let $\alpha_0 = x$, $\beta_0 = y$. Define

$$\alpha_{k+1} = \alpha_k * \beta_k, \quad \beta_{k+1} = \beta_k * \alpha_k.$$

Then α_n and β_n have the same signature up to level n .

e.g.: $\alpha_3 = xyyxyxy$ and $\beta_3 = yxxyxyyx$ are two lattice paths with 8 steps, and have the same signatures up to level three.

$\text{Sig}^{(n)}$ describes finer (local) information of the path when n becomes bigger.

Uniqueness

Chen (1950's): **continuously differentiable curves** are determined by their signatures.

Hambly-Lyons: **paths of finite lengths** are uniquely determined by their signatures **up to tree-like equivalence**.

If α, β are two paths of finite lengths, then $\text{Sig}(\alpha) = \text{Sig}(\beta)$ if and only if $\alpha * \beta^{-1}$ is **equivalent** to a null path.

Finite length paths can have very subtle tree-cancellations, while there is no such nontrivial equivalence if the curve is continuously differentiable (when parametrised at unit speed).

Boedihardjo-Geng-Lyons-Yang: uniqueness for **rough paths**.

Question: how to reconstruct the reduced path from its signature?

Uniqueness and reconstruction

Uniqueness with semi-constructive proofs:

- [Le Jan-Qian](#): Brownian motion sample paths.
- [Boedihardjo-Geng](#): more general Gaussian processes.
- [Geng](#): deterministic rough paths.

Inversion for axis paths

These are paths that move parallel to Euclidean axes. They have the form

$$\gamma = r_1 e_{i_1} * r_2 e_{i_2} * \cdots * r_N e_{i_N} .$$

Information to recover: **ordered directions** $(e_{i_1}, \dots, e_{i_N})$ and **length of each step** (r_1, \dots, r_N) .

Observation:

- $w = e_{i_1} e_{i_2} \cdots e_{i_N}$ is **square-free** (no two adjacent letters are the same), and $C(w) = r_1 r_2 \cdots r_{i_N} \neq 0$.
- If w' is any other square-free word with length $|w'| \geq N$, then $C(w') = 0$.

Conclusion: there exists a **unique longest square free word** w such that $C(w) \neq 0$, then this word tells the **ordered directions** of the path movement.

Inversion for axis paths

These are paths that move parallel to Euclidean axes. They have the form

$$\gamma = r_1 e_{i_1} * r_2 e_{i_2} * \cdots * r_N e_{i_N} .$$

To recover the lengths, let $w = e_{i_1} \cdots e_{i_N}$ be the **unique longest square-free word** as above, and define

$$w_k := e_{i_1} \cdots e_{i_k}^2 \cdots e_{i_N} .$$

Then

$$C(w_k) = \frac{1}{2} r_1 \cdots r_k^2 \cdots r_N \quad \Rightarrow \quad r_k = \frac{2C(w_k)}{C(w)} .$$

Inversion for axis paths

These are paths that move parallel to Euclidean axes. They have the form

$$\gamma = r_1 e_{i_1} * r_2 e_{i_2} * \cdots * r_N e_{i_N} .$$

- 1 Find the **unique longest square-free word** with non-zero coefficient. This word tells the **ordered directions** of the path movement.
- 2 **Move one level up** and compare the coefficients to recover the **length of each step**.

Rely on special structures of the lattice.

Pfeffer-Seigal-Sturmfels: reconstruct paths that arise from a *fixed dictionary*.

Main reconstruction theorem

Theorem (Lyons, X.)

For every k , by using $\text{Sig}(\gamma)$ up to level $N = \mathcal{O}(k^3 \log k)$, we explicitly construct a piecewise linear path $\tilde{\gamma}$ with k pieces such that

$$\sup_{u \in [0,1]} |\tilde{\gamma}'_u - \gamma'_u| < \varepsilon_k$$

when both are parametrized at unit speed (with respect to ℓ^1 norm), and $\varepsilon_k \rightarrow 0$ as $k \rightarrow +\infty$, the speed depending on modulus of continuity of γ' .

The error $\varepsilon_k = \mathcal{O}(k^{-\frac{\alpha^2}{2}})$ if $\gamma \in \mathcal{C}^{1,\alpha}$.

Commutative in smaller scales; noncommutative in larger scales.

Key: how to rule out noncommutativity in small scales?

Main reconstruction theorem

We work with \mathbf{R}^2 for notational simplicity. The piecewise linear path $\tilde{\gamma}$ has the form

$$\tilde{\gamma} = \tilde{\gamma}_1 * \cdots * \tilde{\gamma}_k,$$

where

$$\tilde{\gamma}_j = \frac{\tilde{L}}{k} \left(a_x^{(j)} \rho_j x + a_y^{(j)} (1 - \rho_j) y \right).$$

Hope: each $\tilde{\gamma}_j$ approximates $\gamma_{[\frac{j-1}{k}, \frac{j}{k}]}$ in the ℓ^1 sense.

- $\rho_j, 1 - \rho_j \in [0, 1]$ represents the **unsigned direction**;
- $a_x^{(j)}, a_y^{(j)} \in \{\pm 1\}$ represents the **sign**;
- $\tilde{L} > 0$ approximates the **ℓ^1 length**.

Recovering the increment

Symmetrisation averages out the order.

Summing over all words of length n with k x 's and $n - k$ y 's:

$$\mathcal{S}(k, n - k) = n! \sum_{w \in \mathcal{W}_{k, n-k}} C(w) = \binom{n}{k} (\Delta x)^k (\Delta y)^{n-k}.$$

Maximum: $\frac{k^*}{n-k^*} \approx \frac{|\Delta x|}{|\Delta y|} \Rightarrow$ recovers **unsigned direction**.

More **robust way** of doing it: find k^* such that

$$\sum_{k: |\frac{k}{n} - \frac{k^*}{n}| < \varepsilon} |\mathcal{S}(k, n - k)| \approx \sum_k |\mathcal{S}(k, n - k)|$$

There are more than one such k^* , but all of them are close to each other.

Move one level up: comparing $\mathcal{S}(k^* + 1, n - k^*)$ and $\mathcal{S}(k^*, n - k^*)$ gives the sign of the x direction.

Symmetrisation

Symmetrising k blocks with block size $2n$:

$$\underbrace{*****}_{2n} e_{i_1} \underbrace{*****}_{2n} e_{i_2} \cdots \cdots e_{i_{k-1}} \underbrace{*****}_{2n} .$$

Key: pattern in block j are roughly determined by $\gamma_{[\frac{j-1}{k}, \frac{j}{k}]}$.

Steps:

- 1 Recovering the **unsigned directions** by **checking non-degeneracy**.
- 2 Recovering the **signs** by **moving one level up**.
- 3 Recovering the **length** by a **scaling argument**.

Remark: only uses level $2nk + k - 1$ and $2nk + k$.

Probabilistic interpretation

Terry told me the following probabilistic interpretation of the signature during my PhD.

Suppose $\gamma : [0, 1] \rightarrow \mathbf{R}^2$ is **monotone** in the sense that $x'_t \geq 0$ and $y'_t \geq 0$ for all $t \in [0, 1]$.

Think of the following **Poisson process** $(\mathcal{X}_t, \mathcal{Y}_t)_{t \in [0, 1]}$:

- \mathcal{X}_t generates letter x with intensity x'_t ; \mathcal{Y}_t generates letter y with intensity y'_t ; **simultaneously and independently**.
- We arrange the letters in the order of their arrival time (up to time 1), getting a **(random) word** \mathcal{W} .
- For example, if there are 5 arrivals in total in $[0, 1]$, say x, y, y, x, y at times $0 \leq u_1 < u_2 < u_3 < u_4 < u_5 \leq 1$, then $\mathcal{W} = xyyxy$.

Probabilistic interpretation of signature:

$$C_\gamma(w) = e^L \Pr(\mathcal{W} = w), \quad L = \ell^1 \text{ length of } \gamma.$$

Probabilistic interpretation

For monotone paths:

$$C_\gamma(w) = e^L \Pr(\mathcal{W} = w)$$

Chang-Duffield-Ni-X.: inversion for monotone paths.

General non-monotone paths? x'_t and y'_t can **change signs**.

- Poisson process $(\mathcal{X}_t, \mathcal{Y}_t)$ with intensities $|x'_t|$ and $|y'_t|$.
- Each letter of arrival also carries a sign: if x arrives at time u , then $+1$ if $x'_u > 0$, and -1 if $x'_u < 0$. Same for y .
- Same random word \mathcal{W} as before, but \mathcal{W} also have a sign — the product of the signs of its letters.

Now, we have

$$C_\gamma(w) = e^L \mathbf{E}[\text{sign}(\mathcal{W}) \cdot \mathbf{1}_{\mathcal{W}=w}] , \quad L = \ell^1 \text{ length of } \gamma .$$

Summary

Consequences of the reconstruction:

- 1 Tail signatures already determine C^1 paths.
- 2 'Verification' that higher level signatures describe finer structures of the path.

Quantitative description? Relevant lower bounds (for large n):
Hambly-Lyons, Boedihardjo-Geng

A reverse question: does a version of Bernstein's theorem hold?

What have we learned?

- 1 Symmetrisation counts the frequency but neglects the order; so it gives local increments.
- 2 A certain non-degeneracy criterion is often needed in recovering the directions (Le Jan-Qian, Boedihardjo-Geng, Geng).

Some questions

- 1 Improve efficiency?

Insertion algorithm by Chang-Lyons;

Algebraic structures explored in Améndola-Friz-Sturmfels,
Pfeffer-Seigal-Sturmfels.

- 2 Inversion for rough paths? (Geng)

- 3 Identify the image of the signatures in the tensor algebra.

Expect to involve highly nontrivial interplays between algebraic structure (group-like) and analytic properties (decay) \rightsquigarrow no clue at this moment.

More reasonable to start with monotone paths first.